1 Overview and Motivation

Thus far, in controls, we have introduced linear models (Note 9), discussed how to learn them from data (Note 10), calculated the asymptotical behavior of the state (Note 11), and developed a way to reach a given target state from a given initial state (Note 12). The next step in our journey is to find a way to reach a given target state from a given initial state in the best way (for some reasonable definition of "best"). We will do this in Note 16, but first we will need to introduce some mathematical tools and ideas which are very powerful in their own right.

First among these tools is the idea of orthonormalization.

Key Idea 1 (Orthonormalization)
Orthonormalization is the technique of making a set of vectors easier to work with. Given a set of vectors, it allows us to generate a set of orthogonal and normalized vectors which has the same span.

In this note, we will cover what orthonormal vectors are, the idea of projection and its connections to orthonormality, the process of orthonormalization via the Gram-Schmidt algorithm, some consequences of orthonormalization, and some applications to speeding up system identification via the QR decomposition.

2 Orthogonality and Orthonormality

Recall that in EECS 16A we defined two vectors as orthogonal if they are at 90° to each other, or equivalently if their dot product is 0. This can be extended to inner products $\langle \cdot, \cdot \rangle$, which are themselves generalizations of dot products, defined in EECS 16A Note 21.

Definition 2 (Orthogonal Vectors)

- Let $\vec{x}, \vec{y} \in \mathbb{R}^n$. Then $\vec{x}$ and $\vec{y}$ are orthogonal if and only if $\langle \vec{x}, \vec{y} \rangle = 0$. (1)
- Let $S$ be a set of vectors. Then $S$ is an orthogonal set if and only if every pair of distinct vectors in $S$ is orthogonal, i.e.,

$$\langle \vec{x}, \vec{y} \rangle = 0 \quad \text{for all } \vec{x}, \vec{y} \in S \text{ with } \vec{x} \neq \vec{y}.$$ (2)

Orthogonality is a very restrictive condition, in the sense that orthogonal sets are necessarily very structured. For example, we have the following result:

Theorem 3
Suppose $S \subseteq \mathbb{R}^n$ is an orthogonal set of nonzero vectors. Then $S$ is linearly independent.

Concept Check: Prove Theorem 3.

Theorem 3 says that orthogonality is a stricter condition than linear independence; every orthogonal set is linearly independent, but also has some additional structure that makes the inner products 0.
Corollary 4
Suppose $S \subseteq \mathbb{R}^n$ is an orthogonal set of nonzero vectors. Then $S$ contains at most $n$ vectors.

This is true because linearly independent sets in $\mathbb{R}^n$ have at most $n$ vectors.

In addition, we may also talk about two sets of vectors being orthogonal to each other.

Definition 5 (Orthogonality of Sets)
Let $S_1, S_2 \subseteq \mathbb{R}^n$ be two sets of vectors. Then $S_1$ and $S_2$ are orthogonal to each other if and only if, for all $\vec{x} \in S_1$ and $\vec{y} \in S_2$, we have that $\vec{x}$ and $\vec{y}$ are orthogonal.

NOTE: We can say one set is orthogonal – this means that the vectors in the set are pairwise orthogonal, in the sense of Definition 2. Or we can say two sets are orthogonal to one another – this means that the vectors in the first set are orthogonal to the vectors in the second set, in the sense of Definition 5. We will keep the meaning clear in these notes, but usually the two different notions are distinguished from context.

For example, if the two sets $S_1$ and $S_2$ are subspaces, then by saying $S_1$ and $S_2$ are orthogonal subspaces we mean that every vector in $S_1$ is orthogonal to every vector in $S_2$ (i.e., Definition 5). This is because non-trivial subspaces have infinitely many vectors, and thus by Corollary 4, they cannot be orthogonal sets in the sense of Definition 2.

Definition 5 is usually useful when $S_1$ and $S_2$ are subspaces. In this context, we have the following result, which will be useful in later notes.

Proposition 6
Let $B_1$ be a basis for the subspace $S_1$, and $B_2$ be a basis for the subspace $S_2$. Then $S_1$ and $S_2$ are orthogonal to each other if and only if $B_1$ and $B_2$ are orthogonal to each other.

Concept Check: Prove Proposition 6.

In other words, Proposition 6 says that to check orthogonality of subspaces, it is only required to check that their respective bases are orthogonal.

We may also talk about a vector $\vec{x}$ being orthogonal to a set $S$, which is just shorthand for saying that the sets $\{\vec{x}\}$ and $S$ are orthogonal.

There is also a second part to the definition of orthonormality – that is, normality.

Definition 7 (Normalized Vectors)
Let $\vec{x} \in \mathbb{R}^n$. Then $\vec{x}$ is normalized (i.e., has unit norm) if and only if $\|\vec{x}\| = 1$.

Now we can define orthonormality.
**Definition 8 (Orthonormal Vectors)**

- Let \( \vec{x}, \vec{y} \in \mathbb{R}^n \). Then \( \vec{x} \) and \( \vec{y} \) are orthonormal if and only if:
  - \( \vec{x} \) and \( \vec{y} \) are orthogonal; and
  - both \( \vec{x} \) and \( \vec{y} \) are unit norm.

- Let \( S \subseteq \mathbb{R}^n \). Then \( S \) is an orthonormal set if and only if \( S \) is an orthogonal set and all vectors in \( S \) have unit norm.

The second definition is equivalent to saying that all vectors in \( S \) are unit norm, and also are pairwise orthogonal.

Sets of orthonormal vectors have large amounts of structure, and they are particularly nice to work with. Along with inheriting all properties of orthogonal sets like Theorem 3, i.e.,

![Diagram showing overlapping circles labeled linearly independent, orthogonal, and orthonormal](image)

We also have properties that are specific to orthonormality, such as the following result, which also provides a way to check for orthonormality.

**Theorem 9**

Let \( S \) be a set of vectors. Then \( S \) is an orthonormal set if and only if

\[
\langle \vec{x}, \vec{y} \rangle = \begin{cases} 
1 & \text{if } \vec{x} = \vec{y} \\
0 & \text{if } \vec{x} \neq \vec{y} 
\end{cases}
\quad \text{for all } \vec{x}, \vec{y} \in S.
\] (3)

**Proof.** Indeed, \( \vec{x} \) and \( \vec{y} \) are orthogonal if and only if \( \langle \vec{x}, \vec{y} \rangle = 0 \). And \( \langle \vec{x}, \vec{x} \rangle = \|\vec{x}\|^2 = 1 \) if and only if \( \vec{x} \) is normalized. Applying this to all \( \vec{x}, \vec{y} \in S \) obtains the claim. \( \square \)

**NOTE:** This gives a way to test whether a given set of vectors is orthonormal. If we have an expression for each vector, we can take pairwise inner products and see that they equal to 1 or 0 depending on the vectors we take the inner product of.

Some orthonormality properties manifest nicely as matrix equations with special matrices. For this reason, we define the classes of matrices whose rows or columns are orthonormal sets.
**Definition 10 (Orthonormal Matrices)**

- A square matrix $Q \in \mathbb{R}^{n \times n}$ whose columns or rows form an orthonormal set is called an **orthonormal matrix**.\(^a\)
- A tall matrix $Q \in \mathbb{R}^{m \times n}$ with $m \geq n$ whose columns form an orthonormal set is said to have **orthonormal columns**.
- A wide matrix $Q \in \mathbb{R}^{m \times n}$ with $m \leq n$ whose rows form an orthonormal set is said to have **orthonormal rows**.

\(^a\)Mathematics literature has some unfortunate terminology for this; namely, in mathematics, such matrices are called **orthogonal**, even if the rows or columns are normalized.

**NOTE:** By Corollary 4, we cannot have a tall matrix with orthonormal rows, or a wide matrix with orthonormal columns.

We now present the most foundational property of matrices with orthonormal rows or columns.

**Theorem 11**

(i) If $Q \in \mathbb{R}^{m \times n}$ with $m \geq n$ is a tall matrix, then $Q$ has orthonormal columns if and only if $Q^\top Q = I_n$.

(ii) If $Q \in \mathbb{R}^{m \times n}$ with $m \leq n$ is a wide matrix, then $Q$ has orthonormal rows if and only if $QQ^\top = I_m$.

(iii) If $Q \in \mathbb{R}^{n \times n}$ is a square matrix, then $Q$ is orthonormal if and only if $Q^\top = Q^{-1}$.

The proof of Theorem 11 is on the longer side and may distract from the overall flow of this note, so it is left to Appendix A.1. We fully expect you to read the proof and understand it. It is completely in-scope for the course.

The property $Q^\top = Q^{-1}$ is possibly the most important property of orthonormal matrices.

**Corollary 12**

If $Q \in \mathbb{R}^{n \times n}$ is a square matrix, then it is orthonormal if and only if it has both orthonormal rows and orthonormal columns.

## 3 Projections

Now we will review the concept of projections, which is very much intertwined with the idea of orthonormality. Recall the following definition from EECS 16A Note 21.

**Definition 13 (Projection)**

Let $\vec{x} \in \mathbb{R}^n$ and define $S \subseteq \mathbb{R}^n$ to be a subspace. Then we define the **projection of $\vec{x}$ onto $S$** to be the point

$$\text{proj}_S(\vec{x}) := \arg\min_{\vec{y} \in S} \|\vec{x} - \vec{y}\|.$$  \(4\)

In other words, $\text{proj}_S(\vec{x})$ is the closest point in $S$ to $\vec{x}$.

Now recall from EECS 16A Note 23 the following property of least squares.
**Theorem 14** (Least Squares is Projection)

Let $A \in \mathbb{R}^{m \times n}$ with $m \geq n$ have full column rank, and $\mathbf{y} \in \mathbb{R}^m$. Then

$$\text{proj}_{\text{Col}(A)}(\mathbf{y}) = A(A^\top A)^{-1}A^\top \mathbf{y}. \quad (5)$$

We can now augment this theorem with our knowledge of orthonormal matrices.

**Theorem 15** (Least Squares with Orthonormal Columns)

Let $Q \in \mathbb{R}^{m \times n}$ with $m \geq n$ have orthonormal columns $\mathbf{q}_1, \ldots, \mathbf{q}_n \in \mathbb{R}^m$, and $\mathbf{y} \in \mathbb{R}^m$. Then

$$\text{proj}_{\text{Col}(Q)}(\mathbf{y}) = QQ^\top \mathbf{y} = \sum_{i=1}^{n} \langle \mathbf{y}, \mathbf{q}_i \rangle \mathbf{q}_i. \quad (6)$$

The proof of Theorem 15 is on the longer side and may distract from the overall flow of this note, so it is left to Appendix A.2. We fully expect you to read the proof and understand it. It is completely in-scope for the course.

The following fact will be useful when doing our own orthonormalization computations.

**Corollary 16** (Least Squares with Orthogonal Columns)

Let $Z \in \mathbb{R}^{m \times n}$ with $m \geq n$ have orthogonal and nonzero columns $\mathbf{z}_1, \ldots, \mathbf{z}_n \in \mathbb{R}^m$. Then

$$\text{proj}_{\text{Col}(Z)}(\mathbf{y}) = \sum_{i=1}^{n} \frac{\langle \mathbf{y}, \mathbf{z}_i \rangle}{\| \mathbf{z}_i \|^2} \mathbf{z}_i = \sum_{i=1}^{n} \frac{\langle \mathbf{y}, \mathbf{z}_i \rangle}{\langle \mathbf{z}_i, \mathbf{z}_i \rangle} \mathbf{z}_i. \quad (7)$$

**Proof.** Since $\{\mathbf{z}_1, \ldots, \mathbf{z}_n\}$ is an orthogonal set of nonzero vectors, $\{\frac{\mathbf{z}_1}{\| \mathbf{z}_1 \|}, \ldots, \frac{\mathbf{z}_n}{\| \mathbf{z}_n \|}\}$ is an orthonormal set with the same span, which is still $\text{Col}(Z)$. Starting with Theorem 15, we have

$$\text{proj}_{\text{Col}(Z)}(\mathbf{y}) = \sum_{i=1}^{n} \langle \mathbf{y}, \mathbf{z}_i \rangle \frac{\mathbf{z}_i}{\| \mathbf{z}_i \|} = \sum_{i=1}^{n} \frac{\langle \mathbf{y}, \mathbf{z}_i \rangle}{\| \mathbf{z}_i \|^2} \mathbf{z}_i = \sum_{i=1}^{n} \frac{\langle \mathbf{y}, \mathbf{z}_i \rangle}{\langle \mathbf{z}_i, \mathbf{z}_i \rangle} \mathbf{z}_i. \quad (8)$$

We can, in fact, abstract away the idea of projecting onto column spaces and instead discuss projections onto arbitrary subspaces.

**Theorem 17** (Projection with Orthogonal or Orthonormal Columns)

Let $S \subseteq \mathbb{R}^m$ be a subspace and $\{\mathbf{s}_1, \ldots, \mathbf{s}_n\}$ be a basis of orthogonal vectors for this subspace. Then

$$\text{proj}_S(\mathbf{x}) = \sum_{i=1}^{n} \frac{\langle \mathbf{x}, \mathbf{s}_i \rangle}{\| \mathbf{s}_i \|^2} \mathbf{s}_i = \sum_{i=1}^{n} \frac{\langle \mathbf{x}, \mathbf{s}_i \rangle}{\langle \mathbf{s}_i, \mathbf{s}_i \rangle} \mathbf{s}_i. \quad (9)$$

If $\{\mathbf{s}_1, \ldots, \mathbf{s}_n\}$ is an orthonormal set, then

$$\text{proj}_S(\mathbf{x}) = \sum_{i=1}^{n} \langle \mathbf{x}, \mathbf{s}_i \rangle \mathbf{s}_i. \quad (10)$$

**Concept Check:** Use Theorem 15 and Corollary 16 to prove Theorem 17.
4 Gram-Schmidt Orthonormalization

Now that we have defined orthonormal vectors and sets, and shown off some of their properties, we can now learn how to convert sets of vectors into orthonormal sets of vectors which have the same span. The easiest way to do this is, for each vector in turn, we subtract off the component of this vector which is in the direction of all the previous vectors we processed. Mathematically, this component is the projection of the vector onto the span of the previous vectors.

So really, we can execute the following pseudocode, given some vectors \( \vec{a}_1, \ldots, \vec{a}_\ell \) that we want to orthonormalize:

- Set \( \vec{q}_1 := \frac{\vec{a}_1}{\|\vec{a}_1\|} \).
- For \( i = 2, 3, \ldots, \ell \):
  - Set \( \vec{z}_i := \vec{a}_i - \text{proj}_{\text{Span}(\vec{q}_1, \ldots, \vec{q}_{i-1})}(\vec{a}_i) \).
  - Set \( \vec{q}_i := \frac{\vec{z}_i}{\|\vec{z}_i\|} \).

Intuitively, this gives us a nice picture (in the case \( \ell = n = 2 \)):

\[ \vec{a}_1 \rightarrow \vec{q}_1 \]
\[ \vec{a}_2 \rightarrow \vec{q}_2 \]
\[ \vec{a}_2 - \text{proj}_{\text{Span}(\vec{q}_1)}(\vec{a}_2) \rightarrow \vec{q}_2 \]
\[ \vec{a}_2 - \text{proj}_{\text{Span}(\vec{q}_1, \vec{q}_2)}(\vec{a}_2) \rightarrow \vec{q}_2 \]

\[ \vec{a}_1 \rightarrow \vec{q}_1 \]
\[ \vec{a}_2 \rightarrow \vec{q}_2 \]
\[ \vec{a}_2 - \text{proj}_{\text{Span}(\vec{q}_1)}(\vec{a}_2) \rightarrow \vec{q}_2 \]
\[ \vec{a}_2 - \text{proj}_{\text{Span}(\vec{q}_1, \vec{q}_2)}(\vec{a}_2) \rightarrow \vec{q}_2 \]

\[ \vec{a}_1 \rightarrow \vec{q}_1 \]
\[ \vec{a}_2 \rightarrow \vec{q}_2 \]
\[ \vec{a}_2 - \text{proj}_{\text{Span}(\vec{q}_1)}(\vec{a}_2) \rightarrow \vec{q}_2 \]
\[ \vec{a}_2 - \text{proj}_{\text{Span}(\vec{q}_1, \vec{q}_2)}(\vec{a}_2) \rightarrow \vec{q}_2 \]

**Figure 1:** Orthonormalization in action, going left to right then top to bottom.

Wikipedia also has a very nice animation of this process.

There are two questions remaining:

- How we would calculate \( \text{proj}_{\text{Span}(\vec{q}_1, \ldots, \vec{q}_{i-1})}(\vec{a}_k) \). Indeed, Theorem 15 tells us that

\[ \text{proj}_{\text{Span}(\vec{q}_1, \ldots, \vec{q}_{i-1})}(\vec{a}_k) = \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{q}_k \rangle \vec{q}_k. \]

(11)

- What to do if any \( \vec{z}_i = \vec{0} \) (and so we cannot divide by its norm to get \( \vec{q}_i \)). In this case it depends what exactly we want to do with the algorithm, but usually we can just throw it out, as it does not contribute to the span.
All of this put together exactly forms the Gram-Schmidt orthogonalization algorithm.

Algorithm 18 Gram-Schmidt Orthogonalization Algorithm

**Input:** A set of vectors \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \subseteq \mathbb{R}^n \).

**Output:** An orthonormal set of vectors \( \{ \vec{q}_1, \ldots, \vec{q}_l \} \) which spans the same set.

**function** GRAMSCHMIDT( \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \) )

\[
\text{for } i = 1, \ldots, \ell \text{ do }
\]
\[
\vec{z}_i := \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k \qquad \triangleright \text{If } i = 1 \text{ then there are no } k \text{ to sum over, so } \vec{z}_1 = \vec{a}_1.
\]
\[
\text{if } \vec{z}_i = \vec{0}_n \text{ then }
\]
\[
\vec{p}_i := \vec{0}_n
\]
\[
\text{else }
\]
\[
\vec{p}_i := \frac{\vec{z}_i}{\| \vec{z}_i \|}
\]
\[
\text{end if}
\]
\[
\text{end for}
\]
\[
\text{return } \{ \vec{q}_1, \ldots, \vec{q}_l \} := \{ \vec{p}_1, \ldots, \vec{p}_\ell \} \text{ with the } \ell - l \text{ zero vectors discarded.}
\]

Here we are being very pedantic with the checks that \( \vec{z}_i = \vec{0}_n \), etc, because this function will be extremely critical in a variety of procedures we will come up with. As such, we would like to have well-defined behavior in every possible case.

The main result related to Algorithm 18 is the following theorem, which states that the Gram-Schmidt process accomplishes what we want to do. This type of theorem is called a *proof of correctness* in the literature.

**Theorem 19** (Correctness of Gram-Schmidt Orthogonalization Algorithm)

Let \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \subseteq \mathbb{R}^n \) be a set of vectors and \( \{ \vec{q}_1, \ldots, \vec{q}_l \} := \text{GRAMSCHMIDT}(\{ \vec{a}_1, \ldots, \vec{a}_\ell \}) \). Then

(i) \( \{ \vec{q}_1, \ldots, \vec{q}_l \} \) is an orthonormal set of vectors.

(ii) \( \text{Span}(\vec{q}_1, \ldots, \vec{q}_l) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_\ell) \).

Further, if \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \) are linearly independent then

(iii) \( \ell = l \). (That is, no vectors are discarded.)

(iv) \( \text{Span}(\vec{q}_1, \ldots, \vec{q}_l) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) \) for all \( 1 \leq i \leq \ell \).

The proof of Theorem 19 is on the longer side and may distract from the overall flow of this note, so it is left to Appendix A.3. We fully expect you to read the proof and understand it. It is completely in-scope for the course.

An example of this algorithm being run is in Section 8.1.

5 Consequences of Orthonormalization

Once we know that we are able to orthonormalize finite sets of vectors, we have the following abstract result.

**Theorem 20** (Existence of Orthonormal Basis)

Let \( S \subseteq \mathbb{R}^n \) be a subspace. There exists an orthonormal basis for \( S \).

**Proof.** Take any basis for \( S \) and run Gram-Schmidt on it.

© UCB EECS 16B, Spring 2022. All Rights Reserved. This may not be publicly shared without explicit permission.
There is also a constructive way to extend orthonormal bases for subspaces to other subspaces.

**Theorem 21** (Extending an Orthonormal Basis)
Let \( S \subseteq T \subseteq \mathbb{R}^n \) be subspaces. Then for every orthonormal basis \( B_S \) of \( S \), there exists an orthonormal basis \( B_T \) for \( T \) that contains \( B_S \).

**NOTE:** We mean that \( B_T \) contains \( B_S \) in the following sense: each of the \( \dim(S) \) vectors in \( B_S \) is also one of the \( \dim(T) \) vectors in \( B_T \).

**Proof.** Take any basis \( B \) of \( T \) and run Gram-Schmidt on \( B_S \cup B \) to get \( B_T \). Because the orthonormal vectors in \( B_S \) are processed first by Gram-Schmidt, \( B_S \) is unchanged by Gram-Schmidt, so \( B_S \) is contained in \( B_T \). And since \( B \) spans \( T \), we know that \( B_S \cup B \) spans \( T \), so \( B_T \) is a basis for \( T \).

**NOTE:** Here we assume that we are using ordered sets (which can essentially be thought of as lists or arrays) to hold our bases. In this context, \( B_S \cup B \) can be thought of as the concatenation of the lists representing \( B_S \) and \( B \), with the elements of \( B_S \) going first.

Algorithmically, we can write this process in the following way.

```plaintext
Algorithm 22 Basis Extension
Input: Subspaces \( S \subseteq T \subseteq \mathbb{R}^n \), an orthonormal basis \( B_S \) for \( S \).
Output: An orthonormal basis \( B_T \) for \( T \) that contains \( B_S \).

function EXTEND_BASIS(B_S, T)
    B := any basis for T.
    return \( B_T := \text{GRAMSCHMIDT}(B_S \cup B) \).
end function
```

**NOTE:** The basis extension algorithm is most commonly used with \( T = \mathbb{R}^n \), i.e., in the context of extending orthonormal bases for a subspace \( S \) to an orthonormal basis for \( \mathbb{R}^n \). One common choice for an arbitrary basis \( B \) of \( T \) is to pick \( B \) to be the columns of the identity matrix, i.e., \( B = \{ \vec{e}_1, \ldots, \vec{e}_n \} \).

An example of this algorithm being run is in Section 8.2.

### 6 (OPTIONAL) The QR Decomposition

Theorem 19 tells us that, if \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \) is linearly independent, then each \( \vec{a}_i \) can be written in terms of \( \vec{q}_1, \ldots, \vec{q}_i \) only. This gives the linear system writing \( \vec{a}_1, \ldots, \vec{a}_\ell \) in terms of the \( \vec{q}_k \) a triangular structure. More formally, we have the following decomposition:

**Theorem 23** (QR Decomposition)
Let \( A \in \mathbb{R}^{n \times \ell} \) have full column rank. Then there exists a matrix \( Q \in \mathbb{R}^{n \times \ell} \) with orthonormal columns, and an upper triangular matrix \( R \in \mathbb{R}^{\ell \times \ell} \) (i.e., \( r_{ij} = 0 \) for \( i > j \)), such that

\[
A = QR. \quad (12)
\]

The proof of Theorem 23 is on the longer side and may distract from the overall flow of this note, so it is left to Appendix A.4.

The proof also reveals a constructive way to reach a QR decomposition.
Algorithm 24 QR Decomposition

Input: A matrix $A \in \mathbb{R}^{n \times \ell}$ of full column rank.
Output: A matrix $Q$ with orthonormal columns, and an upper-triangular matrix $R$, such that $A = QR$. 

function $QR(A)$ 
\begin{align*}
\{\vec{a}_1, \ldots, \vec{a}_\ell\} &:= \text{columns of } A \\
\{\vec{q}_1, \ldots, \vec{q}_\ell\} &:= \text{GRAMSCHMIDT}(\{\vec{a}_1, \ldots, \vec{a}_\ell\}) \\
Q &:= \begin{bmatrix}
\langle \vec{a}_1, \vec{q}_1 \rangle & \langle \vec{a}_2, \vec{q}_1 \rangle & \cdots & \langle \vec{a}_\ell, \vec{q}_1 \rangle \\
0 & \langle \vec{a}_2, \vec{q}_2 \rangle & \cdots & \langle \vec{a}_\ell, \vec{q}_2 \rangle \\
\vdots & \vdots & \ddots & \vdots \\
0 & 0 & \cdots & \langle \vec{a}_\ell, \vec{q}_\ell \rangle \\
\end{bmatrix} \\
R &:= \begin{bmatrix}
\vec{q}_1 & \vec{q}_2 & \cdots & \vec{q}_\ell \\
\end{bmatrix} \\
\end{align*}
return $Q, R$
end function

7 (OPTIONAL) Application: Speeding Up Model Selection

Previously, in system identification (see Note 10), we knew exactly how many parameters our model had, and were able to design an algorithm to learn the parameters from data. In practice, we sometimes do not know how many parameters our model has; this is related to the problem of model selection in machine learning literature.

More formally, suppose we have several datasets $D_1 \in \mathbb{R}^{n \times 1}$, $D_2 \in \mathbb{R}^{n \times 2}$, ..., $D_\ell \in \mathbb{R}^{n \times \ell}$, where the columns of $D_i$ are a strict subset of the columns of $D_{i+1}$. In the context of system ID, $D_i$ is the dataset if we include the $i$ “most relevant” states and try to build a model off of that.\footnote{There are two questions here. First, why would we not use all the states? The answer is that in most modern applications, there are usually many possible states to include, and most of them don’t have any effect on the results, so including them in the model is both computationally inefficient and can cause an “overfitting” phenomenon where the parameters use the irrelevant states to erroneously capture the effects of noise. Second, how would we pick the “most relevant” states? The answer is not not simple; there are many ways to do this, and it is a very large problem in statistics and control theory, but it suffices to say that there are many approaches to this “feature selection” that work in practice.} We also have an output vector $\vec{s} \in \mathbb{R}^n$. We want to learn parameter vectors $\vec{p}_1, \ldots, \vec{p}_\ell$ where $\vec{p}_i \in \mathbb{R}^i$ such that $D_i \vec{p}_i \approx \vec{s}$. With these estimates, we can evaluate which $\vec{p}_i$ performs best in practice (or via validation), and thereby learn a compact (minimal states used) yet useful model for the system.

One way to do this is our system identification technique of least squares. In particular, we solve a sequence of least squares problems

$$
\vec{p}_i \in \arg\min_{\vec{p}_i \in \mathbb{R}^i} \|D_i \vec{p}_i - \vec{s}\|
$$

which has solution

$$
\vec{p}_i := (D_i^\top D_i)^{-1} D_i^\top \vec{s}.
$$

The issue is that for large $n$, computing $(D_i^\top D_i)^{-1}$ for every $i$ is computationally inefficient, as measured by the number of floating point operations required for an algorithm to do the computation. Matrix multiplication to compute $D_i^\top D_i$ takes roughly $i^2 n^2$ floating point operations, and then taking the inverse takes roughly $i^3$ floating point operations. Summing this up over all $i \in \{1, \ldots, \ell\}$ means that this process takes roughly $\ell^3 + \ell^2 n^2$ floating point operations, which is extremely large (and thus extremely slow to compute) when $n$ and $\ell$ are large.

Instead, what we can do is use the fact that least squares is a projection. The following algorithm speeds up this repeated least squares a lot.

- Take the QR decomposition $D_i$ to get matrices $Q_i := [\vec{q}_1 \ldots \vec{q}_\ell]$ and $R_\ell$ as defined in the QR decomposition, which takes around $\ell^2 n$ floating point operations. In addition, for notation’s sake we can define $Q_i := [\vec{q}_1 \ldots \vec{q}_\ell]$. 

© UCB EECS 16B, Spring 2022. All Rights Reserved. This may not be publicly shared without explicit permission.
• Let $\vec{w}_1 := \vec{q}_1^\top \vec{s}$; this amounts to efficiently calculating $Q_1^\top \vec{s}$.

• For each $i \in \{2, \ldots, \ell\}$, form $\vec{w}_i$ by appending the entry $\vec{q}_i^\top \vec{s}$ to $\vec{w}_{i-1}$; this amounts to efficiently calculating $Q_i^\top \vec{s}$.

• Note that $\vec{w}_i \neq \vec{p}_i$, because we changed the data from $D_i$ to $Q_i$. To relate the two, we note that

$$D_i \vec{p}_i = D_i (D_i^\top D_i)^{-1} D_i^\top \vec{s} = \text{proj}_{\text{Col}(D_i)}(\vec{s}) = Q_i (Q_i^\top \vec{s}) = Q_i \vec{w}_i. \tag{15}$$

In particular, decomposing $D_i = Q_i R_i$ where $R_i$ is the upper left $i \times i$ sub-block of $R_\ell$ (and is in particular upper triangular), we have

$$Q_i R_i \vec{p}_i = Q_i \vec{w}_i. \tag{16}$$

Left-multiplying by $Q_i$ on both sides, we get

$$R_i \vec{p}_i = \vec{w}_i. \tag{17}$$

Since $R_i$ is upper triangular, we can solve for $\vec{p}_i$ efficiently using backward substitution.

Overall this takes much fewer floating point operations – certainly fewer than $\ell^4 + \ell^2 n^2$.

**Concept Check:** How many floating point operations does our new algorithm require, in terms of $\ell$ and $n$?

## 8 Examples

We give an extended example for the two main algorithms we present in this note.

### 8.1 Gram-Schmidt

Say we have two vectors in $\mathbb{R}^3$:

$$\vec{a}_1 := \begin{bmatrix} 3 \\ 4 \\ 0 \end{bmatrix}, \quad \vec{a}_2 := \begin{bmatrix} 4 \\ 3 \\ 0 \end{bmatrix} \tag{18}$$

that we want to run Gram-Schmidt on. Mechanically, we have

$$\vec{p}_1 := \frac{\vec{a}_1}{\|\vec{a}_1\|} = \frac{1}{5} \begin{bmatrix} 3 \\ 4 \\ 0 \end{bmatrix} = \frac{3}{5} \begin{bmatrix} 1 \\ 4/3 \\ 0 \end{bmatrix}. \tag{19}$$

Now we compute the projection residual:

$$\vec{z}_2 = \vec{a}_2 - \langle \vec{a}_2, \vec{p}_1 \rangle \vec{p}_1 \tag{20}$$

$$= \begin{bmatrix} 4 \\ 3 \\ 0 \end{bmatrix} - \left( \begin{bmatrix} 4 \\ 3 \\ 0 \end{bmatrix} \cdot \frac{3}{5} \right) \begin{bmatrix} 3 \\ 4/3 \\ 0 \end{bmatrix} \tag{21}$$

$$= \begin{bmatrix} 4 \\ 3 \\ 0 \end{bmatrix} - \left( \frac{4 \cdot 3}{5} + \frac{3 \cdot 4}{5} + 0 \cdot 0 \right) \begin{bmatrix} 3 \\ 4/3 \\ 0 \end{bmatrix} \tag{22}$$

$$= \begin{bmatrix} 4 \\ 3 \\ 0 \end{bmatrix} - \frac{24}{5} \begin{bmatrix} 3 \\ 4/3 \\ 0 \end{bmatrix} \tag{23}$$
Then we can normalize it to get $\vec{p}_2$ as

$$\vec{p}_2 = \frac{\vec{z}_2}{\|\vec{z}_2\|} = \begin{bmatrix} 4 \frac{5}{3} \\ -3 \frac{5}{3} \\ 0 \end{bmatrix}. \quad (27)$$

Since $\vec{p}_1 \neq \vec{0}_3$ and $\vec{p}_2 \neq \vec{0}_3$, we output

$$\vec{q}_1 = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix} \quad \text{and} \quad \vec{q}_2 = \begin{bmatrix} 4 \frac{5}{3} \\ -3 \frac{5}{3} \\ 0 \end{bmatrix}. \quad (28)$$

### 8.2 Extending Basis

Let’s say we wanted to extend our basis $\{\vec{q}_1, \vec{q}_2\}$ to a basis for $\mathbb{R}^3$. To do this, we would pick a basis for $\mathbb{R}^3$, usually the identity basis $\{\vec{e}_1, \vec{e}_2, \vec{e}_3\}$. Then we would run Gram-Schmidt on $\{\vec{q}_1, \vec{q}_2, \vec{e}_1, \vec{e}_2, \vec{e}_3\}$. Since $\mathbb{R}^3$ is three-dimensional, any basis for $\mathbb{R}^3$ contains 3 vectors, so the output of Gram-Schmidt will also contain 3 vectors, which is an orthonormal basis for $\mathbb{R}^3$.

Anyways, we re-run Gram-Schmidt on this basis. The computation for $\vec{p}_1$ and $\vec{p}_2$ goes the same way as the previous section since we are running it on the same vectors in the same order, so we get

$$\vec{p}_1 := \begin{bmatrix} \frac{3}{5} \\ 0 \\ 0 \end{bmatrix} \quad \text{and} \quad \vec{p}_2 := \begin{bmatrix} 4 \frac{5}{3} \\ -3 \frac{5}{3} \\ 0 \end{bmatrix}. \quad (29)$$

Now $\vec{p}_3, \vec{p}_4, \vec{p}_5$ are new. Let’s start with computing the projection residual:

$$\vec{z}_3 = \vec{e}_3 - \langle \vec{e}_3, \vec{p}_1 \rangle \vec{p}_1 - \langle \vec{e}_3, \vec{p}_2 \rangle \vec{p}_2$$

$$= \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \left[ \begin{bmatrix} \frac{3}{5} \\ \frac{4}{5} \\ 0 \end{bmatrix} \right] - \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} \left[ \begin{bmatrix} \frac{4}{5} \\ -\frac{3}{5} \\ 0 \end{bmatrix} \right] \left[ \begin{bmatrix} \frac{4}{5} \\ -\frac{3}{5} \\ 0 \end{bmatrix} \right]$$

$$= \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} \frac{3}{5} \\ \frac{4}{5} \\ 0 \end{bmatrix} - \begin{bmatrix} \frac{4}{5} \\ -\frac{3}{5} \\ 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} 0 \\ \frac{12}{25} \\ 0 \end{bmatrix} - \begin{bmatrix} 16 \\ -12 \\ 0 \end{bmatrix}$$

$$= \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}. \quad (34)$$

This implies that $\vec{p}_3 = \vec{0}_3$ and therefore there is no corresponding $\vec{q}$ vector. Similarly,

$$\vec{z}_4 = \vec{e}_4 - \langle \vec{e}_4, \vec{p}_1 \rangle \vec{p}_1 - \langle \vec{e}_4, \vec{p}_2 \rangle \vec{p}_2 - \langle \vec{e}_4, \vec{p}_3 \rangle \vec{p}_3$$

$$= \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}. \quad (35)$$
\[ \vec{e}_2 - \langle \vec{e}_2, \vec{p}_1 \rangle \vec{p}_1 - \langle \vec{e}_2, \vec{p}_2 \rangle \vec{p}_2 \]

\[ = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} - \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} \frac{3}{4} \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} 0 \\ 0 \\ \frac{4}{5} \end{bmatrix} \begin{bmatrix} \frac{4}{5} \\ -\frac{3}{5} \\ 0 \end{bmatrix} \]

\[ = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \]

\[ = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} - \begin{bmatrix} \frac{3}{4} \\ 0 \\ 0 \end{bmatrix} - \begin{bmatrix} \frac{4}{5} \\ -\frac{3}{5} \\ 0 \end{bmatrix} \]

\[ = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \]

Again, this implies that \( \vec{p}_4 = \vec{0}_3 \) and therefore there is no corresponding \( \vec{q} \) vector.

Finally, computing the residual for \( \vec{e}_3 \) gets us

\[ \vec{z}_5 = \vec{e}_3 - \langle \vec{e}_3, \vec{p}_1 \rangle \vec{p}_1 - \langle \vec{e}_3, \vec{p}_2 \rangle \vec{p}_2 - \langle \vec{e}_3, \vec{p}_3 \rangle \vec{p}_3 - \langle \vec{e}_3, \vec{p}_4 \rangle \vec{p}_4 \]

\[ = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} - \begin{bmatrix} 0 \\ 0 \\ \frac{3}{4} \end{bmatrix} - \begin{bmatrix} \frac{4}{5} \\ -\frac{3}{5} \\ 0 \end{bmatrix} \]

\[ = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \]

Then we have

\[ \vec{p}_5 := \frac{\vec{z}_5}{\|\vec{z}_5\|} = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \]

Thus \( \vec{q}_3 = \vec{p}_5 = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \). Gram-Schmidt outputs

\[ \vec{q}_1 = \begin{bmatrix} \frac{3}{4} \\ 0 \\ \frac{4}{5} \end{bmatrix}, \quad \vec{q}_2 = \begin{bmatrix} \frac{3}{4} \\ 0 \\ -\frac{3}{5} \end{bmatrix}, \quad \vec{q}_3 = \begin{bmatrix} 0 \\ 0 \\ 1 \end{bmatrix} \]

which is indeed an orthonormal basis for \( \mathbb{R}^3 \).
A Longer Proofs

A.1 Proof of Theorem 11

Proof of Theorem 11.

(i) Suppose the columns of $Q$ are $\vec{q}_1, \ldots, \vec{q}_n \in \mathbb{R}^m$. Then we have

$$Q^\top Q = [\vec{q}_1 \cdots \vec{q}_n]^\top [\vec{q}_1 \cdots \vec{q}_n]$$

(48)

$$= \left[\begin{array}{c} \vec{q}_1^\top \\ \vdots \\ \vec{q}_n^\top \end{array}\right] \left[\begin{array}{c} \vec{q}_1 \\ \vdots \\ \vec{q}_n \end{array}\right]$$

(49)

$$= \left[\begin{array}{ccc} \vec{q}_1^\top \vec{q}_1 & \cdots & \vec{q}_1^\top \vec{q}_n \\ \vdots & \ddots & \vdots \\ \vec{q}_n^\top \vec{q}_1 & \cdots & \vec{q}_n^\top \vec{q}_n \end{array}\right]$$

(50)

$$= \left[\begin{array}{cc} \langle \vec{q}_1, \vec{q}_1 \rangle & \cdots & \langle \vec{q}_1, \vec{q}_n \rangle \\ \vdots & \ddots & \vdots \\ \langle \vec{q}_n, \vec{q}_1 \rangle & \cdots & \langle \vec{q}_n, \vec{q}_n \rangle \end{array}\right]$$

(51)

By Theorem 9, we see that $Q^\top Q = I_n$ if and only if $\{\vec{q}_1, \ldots, \vec{q}_n\}$ are orthonormal.

(ii) Apply (i) to $Q^\top \in \mathbb{R}^{n \times m}$, which is a tall matrix. Note that

$$QQ^\top = (Q^\top)^\top (Q^\top).$$

(52)

This tells us that

$$QQ^\top = I_m$$

(53)

if and only if $Q^\top$ has orthonormal columns, i.e., $Q$ has orthonormal rows.

(iii) Suppose $Q$ is orthogonal. Then it has orthonormal columns or orthonormal rows. If it has orthonormal columns then, by (i), $Q^\top Q = I_n$; if it has orthonormal rows then, by (ii), $QQ^\top = I_n$. Either way, $Q^\top = Q^{-1}$ as desired. Going the other way, if $Q^\top = Q^{-1}$ then $Q^\top Q = I_n$, which confirms by (i) that $Q$ has orthonormal columns and is thus orthonormal.

\[\square\]

A.2 Proof of Theorem 15

Proof of Theorem 15. Since $Q$ has orthonormal columns, by Theorem 3, we know that the columns of $Q$ are linearly independent, so $Q$ has full column rank. By Theorem 14, we have

$$\text{proj}_{\text{Col}(Q)}(\vec{y}) = Q(Q^\top Q)^{-1}Q^\top \vec{y}.$$  

(54)

By Theorem 11, since $Q$ has orthonormal columns, we know that $Q^\top Q = I_n$. Thus

$$\text{proj}_{\text{Col}(Q)}(\vec{y}) = Q(I_n)^{-1}Q^\top \vec{y}$$

(55)

$$= Q(I_n)^{-1}Q^\top \vec{y}$$

(56)

$$= QQ^\top \vec{y}$$

(57)
we will start with proving facts about the what we want it to do. Given that the general flow of the algorithm is to build from the usual flow for a proof of correctness is to show that each line, or block of lines, in the algorithm does For notation’s sake, let Let A.3 Proof of Theorem 19

EECS 16B Note 13: Orthonormalization 2022-03-10 22:14:25-08:00

⃗(a) We want to show some claims about the ⃗(II) For each ⃗(I) ⃗(III) ⃗(62) ⃗(63) ⃗(64) ⃗(65) ⃗(66)

\[
\begin{align*}
\langle \vec{y}, \vec{q}_1 \rangle & = \langle \vec{y}, \vec{q}_1 \rangle \\
\langle \vec{y}, \vec{q}_1 \rangle & = \langle \vec{y}, \vec{q}_1 \rangle \\
\langle \vec{y}, \vec{q}_1 \rangle & = \sum_{i=1}^{n} \langle \vec{y}, \vec{q}_i \rangle \vec{q}_i. \\
\end{align*}
\]

\[a_i \in \{a_1, \ldots, a_\ell\}. \]

The usual flow for a proof of correctness is to show that each line, or block of lines, in the algorithm does what we want it to do. Given that the general flow of the algorithm is to build from $\vec{a}_i$'s to $\vec{z}_i$'s to $\vec{p}_i$'s to $\vec{a}_i$'s, we will start with proving facts about the $\vec{z}_i$'s.

(a) We want to show some claims about the $\vec{z}_i$'s:

(I) $\{\vec{z}_1, \ldots, \vec{z}_\ell\}$ is an orthogonal set.

(II) For each $i$ we have $\text{Span}(\vec{z}_1, \ldots, \vec{z}_i) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i)$.

(III) $\vec{z}_i = \vec{0}_n$ if and only if $\vec{a}_i \in \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1})$.

To do this, we first write the $\vec{z}_i$ in terms of $\vec{a}_i$ and other $\vec{z}_k$.

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k = \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \langle \vec{a}_i, \vec{p}_k \rangle \vec{p}_k
\]
This is an expression for \( \vec{z}_i \) in terms of \( \vec{z}_k \) for \( k < i \), as well as \( \vec{a}_i \). Note that we summed over only \( k \) such that \( \vec{z}_k \neq \vec{0}_n \); this is to avoid division by 0 errors.

Now we are ready to prove our claims.

(i) We show that \( \{ \vec{z}_1, \ldots, \vec{z}_i \} \) is an orthogonal set. We do this by showing that for each \( i \), \( \vec{z}_i \) is orthogonal to all the \( \vec{z}_k \) for \( k < i \). We may show this by using a proof that relies on recursion.\(^2\)

Our recursive base case is \( i = 1 \), where there are no \( \vec{z}_j \) for \( k < i \), and so we are done by definition – we have shown that \( \vec{z}_1 \) is orthogonal to all the \( \vec{z}_k \) coming before it (because there are no such \( \vec{z}_k \) and the statement becomes tautological).

For our general recursive case, we show that \( \vec{z}_i \) is orthogonal to all \( \vec{z}_k \) before it, i.e., for \( k < i \). If \( \vec{z}_k = \vec{0}_n \) then

\[
\langle \vec{z}_i, \vec{z}_k \rangle = \langle \vec{z}_i, \vec{0}_n \rangle = 0
\]  

so \( \vec{z}_i \) is orthogonal to \( \vec{z}_k \). We now just need to show that \( \vec{z}_i \) is orthogonal to all \( \vec{z}_k \neq \vec{0}_n \) for \( k < i \). These are exactly the terms appearing in the sum in Equation (67). Indeed, we can take the inner product and use Equation (67) to get

\[
\langle \vec{z}_i, \vec{z}_k \rangle = \left( \sum_{j=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_j \rangle}{\langle \vec{z}_j, \vec{z}_j \rangle} \langle \vec{z}_j, \vec{z}_k \rangle \right)
\]

(68)

Here we used bilinearity of the inner product to separate out terms and move scalars to the outside of the inner product. So now we have a bunch of inner products of the form \( \langle \vec{z}_j, \vec{z}_k \rangle \) for \( 1 \leq j, k < i \). We can now recurse on the smaller problems; by the so-called "recursive leap of faith", we have that \( \vec{z}_j \) is orthogonal to all the \( \vec{z}_m \) for \( m < j \), and similarly that \( \vec{z}_k \) is orthogonal to all the \( \vec{z}_m \) for \( m < k \). Thus, if one of \( j \) or \( k \) is less than the other, then they are orthogonal; otherwise, \( j = k \).

**NOTE:** It is possible to use recursion here since \( j, k < i \) and thus the argument size is decreasing in the "recursive call", so eventually it will decrease to 1 and hit our base case.

Anyways, now that we know \( \langle \vec{z}_j, \vec{z}_k \rangle = 0 \) unless \( j = k \), we can write

\[
\langle \vec{z}_i, \vec{z}_k \rangle = \left( \sum_{j=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_j \rangle}{\langle \vec{z}_j, \vec{z}_j \rangle} \langle \vec{z}_j, \vec{z}_k \rangle \right)
\]

(72)

\[
= \langle \vec{a}_i, \vec{z}_k \rangle - \left( \sum_{j=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_j \rangle}{\langle \vec{z}_j, \vec{z}_j \rangle} \langle \vec{z}_j, \vec{z}_k \rangle \right)
\]

(73)

\[
= \langle \vec{a}_i, \vec{z}_k \rangle - \left( \sum_{j=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_j \rangle}{\langle \vec{z}_j, \vec{z}_j \rangle} \langle \vec{z}_j, \vec{z}_k \rangle \right) = 0
\]

(74)

\[
= \langle \vec{a}_i, \vec{z}_k \rangle - \langle \vec{a}_i, \vec{z}_k \rangle
\]

(75)

\(^2\)If we had CS70 or an equivalent class as a prerequisite, this would be essentially an induction proof. But again, knowing the formalism for induction isn’t needed for reading and understanding this.
Finally, we want to show that \( \vec{z}_i \) is orthogonal to \( \vec{z}_k \). This shows that \( \vec{z}_i \) is orthogonal to \( \vec{z}_1, \ldots, \vec{z}_{i-1} \) for each \( i \). Thus we have that \( \{\vec{z}_1, \ldots, \vec{z}_i\} \) is an orthogonal set for every \( i \), and so \( \{\vec{z}_1, \ldots, \vec{z}_k\} \) is an orthogonal set.

**II** We show that \( \text{Span}(\vec{z}_1, \ldots, \vec{z}_i) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) \) for each \( i \). This is equivalent to showing that \( \vec{z}_i \) can be written as a linear combination of \( \vec{a}_1, \ldots, \vec{a}_i \), and \( \vec{a}_i \) can be written as a linear combination of \( \vec{z}_1, \ldots, \vec{z}_i \), for each \( i \).

We first claim that \( \vec{a}_i \) can be written as a linear combination of \( \vec{z}_1, \ldots, \vec{z}_i \), for each \( i \). Indeed, by rearranging Equation (67), we have

\[
\vec{a}_i = \vec{z}_i + \sum_{k=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_k \rangle}{\langle \vec{z}_k, \vec{z}_k \rangle} \vec{z}_k.
\]  

This is a linear combination of \( \vec{z}_1, \ldots, \vec{z}_i \) whose value is \( \vec{a}_i \), and so we are done.

Now we claim that \( \vec{z}_i \) can be written as a linear combination of \( \vec{a}_1, \ldots, \vec{a}_i \). We do this via a recursive approach again.

The recursive base case is \( i = 1 \), in which case \( \vec{z}_1 = \vec{a}_1 \) is the linear combination we want.

The general recursive case starts again with Equation (67), giving us:

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_k \rangle}{\langle \vec{z}_k, \vec{z}_k \rangle} \vec{z}_k.
\]  

Since \( k < i \), we can recursively write \( \vec{z}_k = \sum_{j=1}^{k} \alpha_{jk} \vec{a}_k \), i.e., generate a linear combination of \( \vec{a}_1, \ldots, \vec{a}_k \) whose value is \( \vec{z}_k \). This gets us

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_k \rangle}{\langle \vec{z}_k, \vec{z}_k \rangle} \vec{z}_k
\]  

which is a linear combination of \( \vec{a}_1, \ldots, \vec{a}_i \) (remember that the term \( \frac{\langle \vec{a}_i, \vec{z}_k \rangle}{\langle \vec{z}_k, \vec{z}_k \rangle} \) is just a scalar). Thus we are done and the spans are equal.

**III** Finally, we want to show that \( \vec{z}_i = \vec{0}_a \) if and only if \( \vec{a}_i \in \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1}) \). Indeed, we start with Equation (67), use Theorem 17 to write the sum as a projection, and then simplify with the fact that adding zero vectors to a set maintains the span of the set.

\[
\vec{z}_i = \vec{a}_i - \sum_{k=1}^{i-1} \frac{\langle \vec{a}_i, \vec{z}_k \rangle}{\langle \vec{z}_k, \vec{z}_k \rangle} \vec{z}_k
\]  

\[
= \vec{a}_i - \text{proj}_{\text{Span}(\vec{z}_1, \ldots, \vec{z}_{i-1})} (\vec{a}_i)
\]  

We can now use (II) to simplify even more, since \( \text{Span}(\vec{z}_1, \ldots, \vec{z}_{i-1}) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1}) \).

\[
\vec{z}_i = \vec{a}_i - \text{proj}_{\text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1})} (\vec{a}_i)
\]  

© UCB EECS 16B, Spring 2022. All Rights Reserved. This may not be publicly shared without explicit permission.
Theorem 19. Suppose for now that

\[ \vec{a}_i = \text{proj}_{\text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1})}(\vec{a}_i). \]  

(86)

Thus \( \vec{z}_i = \vec{0}_n \) if and only if \( \vec{a}_i = \text{proj}_{\text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1})}(\vec{a}_i) \), i.e., if and only if \( \vec{a}_i \in \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1}) \).

Phew! That was quite a lot. What just transpired is one of the single most involved proofs in this course. But it’s necessary to completely understand Gram-Schmidt, which is an extremely critical sub-routine in many concepts we will cover.

If you’ve come this far, take a breather, and then dive right back into it. We promise, the rest of the proof is a lot simpler.

(b) We now want to show some claims about the \( \vec{p}_i \)'s:

(I) \( \{\vec{p}_1, \ldots, \vec{p}_i\} \) is an orthogonal set.

(II) For every \( i \), either \( \vec{p}_i = \vec{0}_n \) or \( \|\vec{p}_i\| = 1 \).

(III) For each \( i \) we have \( \text{Span}(\vec{p}_1, \ldots, \vec{p}_i) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) \).

(IV) \( \vec{p}_i = \vec{0}_n \) if and only if \( \vec{a}_i \in \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1}) \).

We can show them rather quickly, as a result of our previous claims about the \( \vec{z}_i \)'s. The crucial observation to make here is that, for each \( i \), \( \vec{p}_i \) is a scalar multiple of \( \vec{z}_i \), say \( \vec{p}_i = \alpha_i \vec{z}_i \). If \( \vec{z}_i = \vec{0} \) we may take \( \alpha_i = 1 \); otherwise we can take \( \alpha_i = \frac{1}{\|\vec{z}_i\|} \). Thus in either case \( \alpha_i \neq 0 \). With this framework, we now give the proofs.

(I) Let \( 1 \leq k < i \leq \ell \). Then

\[ \langle \vec{p}_i, \vec{p}_k \rangle = \langle \alpha_i \vec{z}_i, \alpha_k \vec{z}_k \rangle = \alpha_i \alpha_k \langle \vec{z}_i, \vec{z}_k \rangle = 0 \]  

(87)

(88)

(89)

since \( \{\vec{z}_1, \ldots, \vec{z}_\ell\} \) is an orthogonal set. Thus \( \{\vec{p}_1, \ldots, \vec{p}_\ell\} \) is an orthogonal set.

(II) Take \( i \) such that \( \vec{z}_i = \vec{0}_n \). Then by the algorithm, \( \vec{p}_i = \vec{0}_n \), so

\[ \|\vec{p}_i\| = \|\vec{0}_n\| = 0. \]  

(90)

Now take \( i \) such that \( \vec{z}_i \neq \vec{0}_n \). Then

\[ \|\vec{p}_i\| = \|\alpha_i \vec{z}_i\| = \alpha_i \|\vec{z}_i\| = \frac{\|\vec{z}_i\|}{\|\vec{z}_i\|} = 1. \]  

(91)

(III) For each \( i \) we have \( \text{Span}(\vec{z}_1, \ldots, \vec{z}_i) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) \). Since \( \vec{p}_i = \alpha_i \vec{z}_i \) where \( \alpha_i \neq 0 \), we know that \( \text{Span}(\vec{z}_1, \ldots, \vec{z}_i) = \text{Span}(\vec{p}_1, \ldots, \vec{p}_i) \). Thus \( \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) = \text{Span}(\vec{p}_1, \ldots, \vec{p}_i) \).

(IV) We have that \( \vec{p}_i = \alpha_i \vec{z}_i \) with \( \alpha_i \neq 0 \), so \( \vec{p}_i = \vec{0}_n \) if and only if \( \vec{z}_i = \vec{0}_n \). We have already shown that this is true if and only if \( \vec{a}_i \in \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1}) \), so we obtain the conclusion.

Nearly there! Once we have the facts proven above, we can prove the required original claims made in Theorem 19. Suppose for now that \( \{\vec{a}_1, \ldots, \vec{a}_\ell\} \) is a general set of vectors.

(i) To show that \( \{\vec{q}_1, \ldots, \vec{q}_\ell\} \) is an orthonormal set, we use the fact that \( \{\vec{p}_1, \ldots, \vec{p}_\ell\} \) is an orthogonal set where each vector is either the zero vector or has unit norm. Since \( \{\vec{a}_1, \ldots, \vec{a}_\ell\} \) is \( \{\vec{p}_1, \ldots, \vec{p}_\ell\} \) with all zero vectors removed, we have that \( \{\vec{q}_1, \ldots, \vec{q}_\ell\} \) is an orthogonal set where all vectors have unit norm, and is hence an orthonormal set.

(ii) To show that \( \text{Span}(\vec{q}_1, \ldots, \vec{q}_\ell) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_\ell) \), note that we already know \( \text{Span}(\vec{p}_1, \ldots, \vec{p}_\ell) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_\ell) \). Since \( \{\vec{q}_1, \ldots, \vec{q}_\ell\} \) is \( \{\vec{p}_1, \ldots, \vec{p}_\ell\} \) with all zero vectors removed, and zero vectors don’t contribute to any span, we have that \( \text{Span}(\vec{q}_1, \ldots, \vec{q}_\ell) = \text{Span}(\vec{p}_1, \ldots, \vec{p}_\ell) \). Putting this together, we have that \( \text{Span}(\vec{q}_1, \ldots, \vec{q}_\ell) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_\ell) \) as desired.
Now suppose that \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \) is a linearly independent set of vectors.

(iii) To show that \( l = \ell \), we show that no \( \vec{p}_i \) is a zero vector that gets discarded. Indeed, we know that \( \vec{p}_i = \vec{0}_n \) if and only if \( \vec{a}_i = \text{Span}(\vec{a}_1, \ldots, \vec{a}_{i-1}) \). But since \( \{ \vec{a}_1, \ldots, \vec{a}_\ell \} \) is a linearly independent set, this never happens, and thus \( \vec{p}_i \neq \vec{0}_n \) for all \( i \). Thus no discards happen, so \( l = \ell \).

(iv) Since \( l = \ell \), we have \( \vec{p}_i = \vec{q}_i \) for all \( i \). Since we know \( \text{Span}(\vec{p}_1, \ldots, \vec{p}_i) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) \) for each \( i \), we also have that \( \text{Span}(\vec{q}_1, \ldots, \vec{q}_i) = \text{Span}(\vec{a}_1, \ldots, \vec{a}_i) \) for each \( i \), as desired.

A.4 Proof of Theorem 23

Proof of Theorem 23. By Theorem 19, we can let \( \{ \vec{q}_1, \ldots, \vec{q}_\ell \} \) be the output of Gram-Schmidt when run on the columns of \( A \) (the theorem invocation is to ensure that we receive \( \ell \) vectors back). we have

\[
\vec{a}_i = \text{proj}_{\text{Span}(\vec{a}_1, \ldots, \vec{a}_i)}(\vec{a}_i) = \text{proj}_{\text{Span}(\vec{q}_1, \ldots, \vec{q}_i)}(\vec{a}_i)
\]

\[
= \sum_{k=1}^{i} \langle \vec{a}_i, \vec{q}_k \rangle \vec{q}_k
\]

\[
= \begin{bmatrix} \vec{q}_1 & \cdots & \vec{q}_i \end{bmatrix} \begin{bmatrix} \langle \vec{a}_i, \vec{q}_1 \rangle \\ \vdots \\ \langle \vec{a}_i, \vec{q}_i \rangle \end{bmatrix}
\]

\[
= \begin{bmatrix} \vec{q}_1 & \cdots & \vec{q}_i & \vec{q}_{i+1} & \cdots & \vec{q}_\ell \end{bmatrix} \begin{bmatrix} \langle \vec{a}_i, \vec{q}_1 \rangle \\ \vdots \\ \langle \vec{a}_i, \vec{q}_i \rangle \\ 0 \\ \vdots \\ 0 \end{bmatrix}
\]

Putting these columns together, we get

\[
A = \begin{bmatrix} \vec{a}_1 & \cdots & \vec{a}_\ell \end{bmatrix} = \begin{bmatrix} \vec{q}_1 & \cdots & \vec{q}_\ell \end{bmatrix} \begin{bmatrix} \langle \vec{a}_1, \vec{q}_1 \rangle & \langle \vec{a}_2, \vec{q}_1 \rangle & \cdots & \langle \vec{a}_\ell, \vec{q}_1 \rangle \\ \langle \vec{a}_1, \vec{q}_2 \rangle & \langle \vec{a}_2, \vec{q}_2 \rangle & \cdots & \langle \vec{a}_\ell, \vec{q}_2 \rangle \\ \vdots & \vdots & \ddots & \vdots \\ \langle \vec{a}_1, \vec{q}_\ell \rangle & \langle \vec{a}_2, \vec{q}_\ell \rangle & \cdots & \langle \vec{a}_\ell, \vec{q}_\ell \rangle \end{bmatrix}
\]

which is the desired decomposition.

Contributors:
- Druv Pai.
- Ashwin Vangipuram.
- Anant Sahai.
- Jennifer Shih.
- Rachel Hochman.
- Vasuki Narasimha Swamy.
- Steven Cao.